

Analyzing and Protecting Communication Metadata

Ludovic Barman

Laboratory for Data Security & Decentralized and Distributed Systems Laboratory

PhD Public Defense, 17.09.2021

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WARNING

Simplified presentation. Some technical details
have been simplified for the sake of clarity.

Please check the last slide for the proper references.

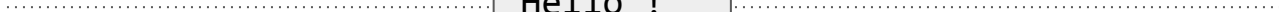
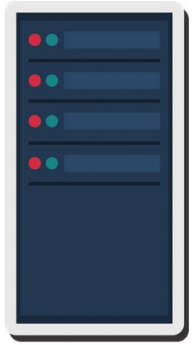






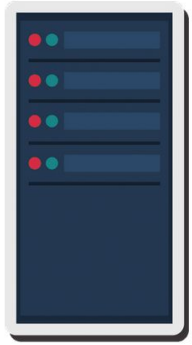






Hello !



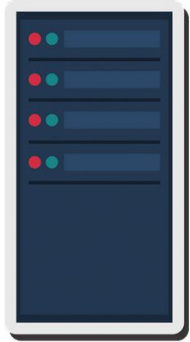


0A6DF31



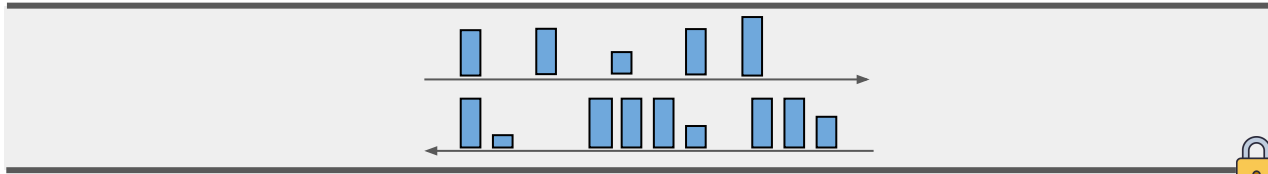
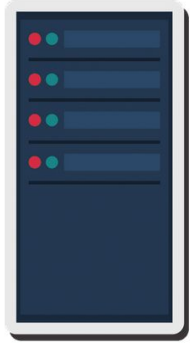






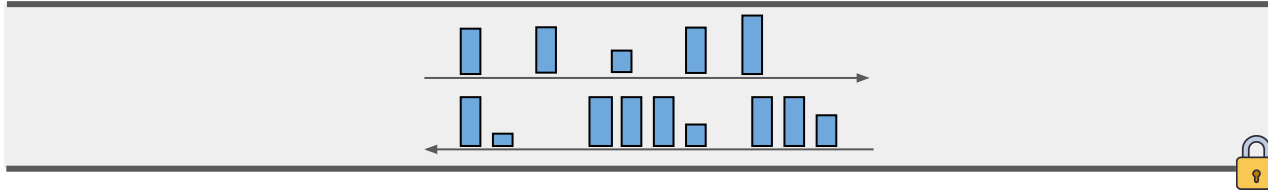
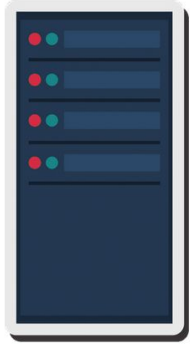
Encrypted connection

implies confidentiality from 3rd parties



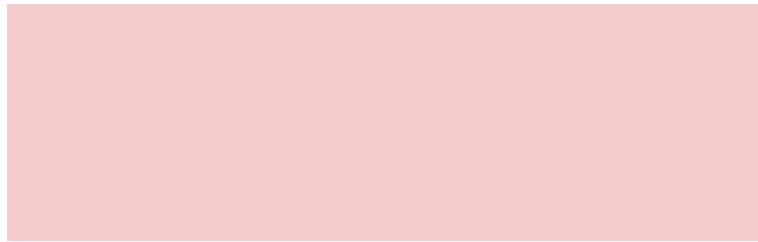
Encrypted connection

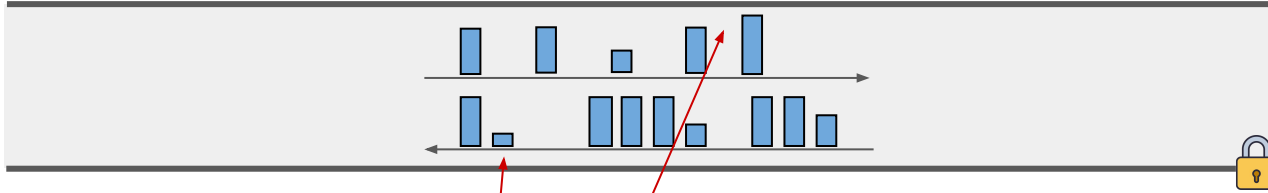
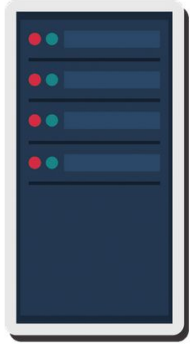
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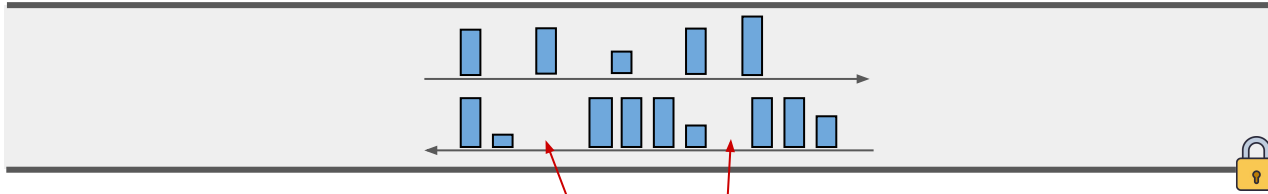
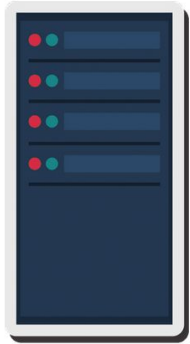




Encrypted connection
implies confidentiality from 3rd parties



Packet sizes



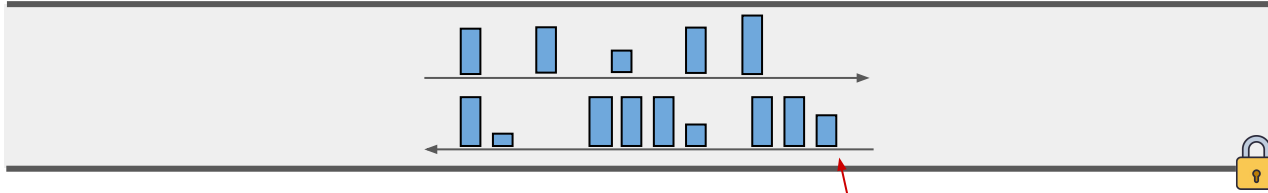
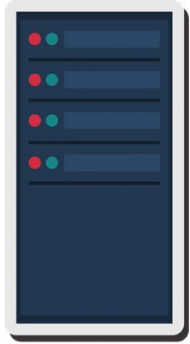
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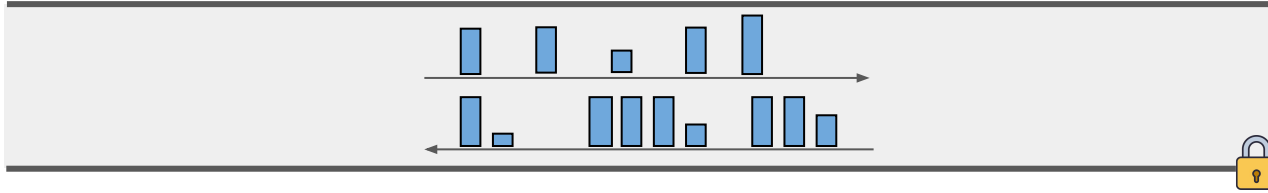
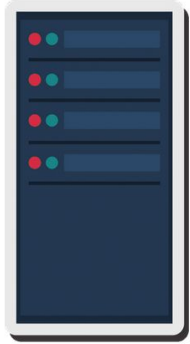
Pauses



Encrypted connection
implies confidentiality from 3rd parties



Packet sizes
Pauses
Total duration



Encrypted connection

implies confidentiality from 3rd parties



Packet sizes

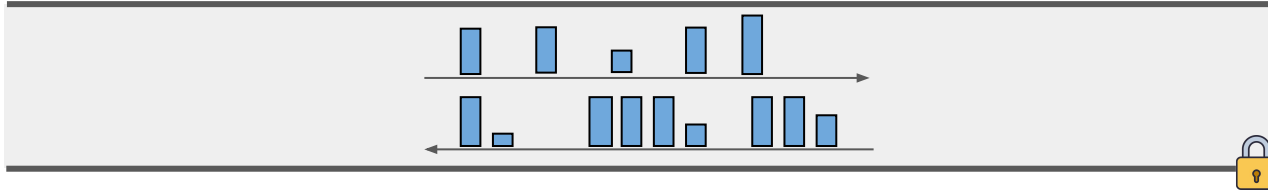
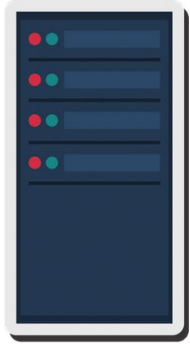
Total duration

Pauses

Packet counts

Ordering

etc.



Encrypted connection

implies confidentiality from 3rd parties



Unprotected *Metadata* :

Packet sizes

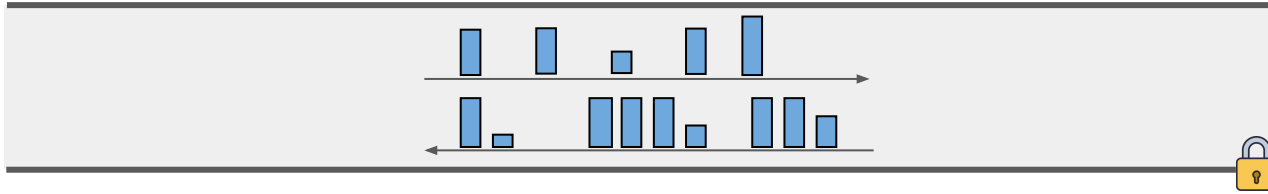
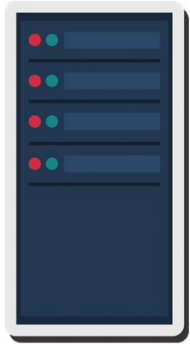
Total duration

Pauses

Packet counts

Ordering

etc.



Encrypted connection

implies confidentiality from 3rd parties



Unprotected *Metadata* :

Packet sizes

Total duration

Pauses

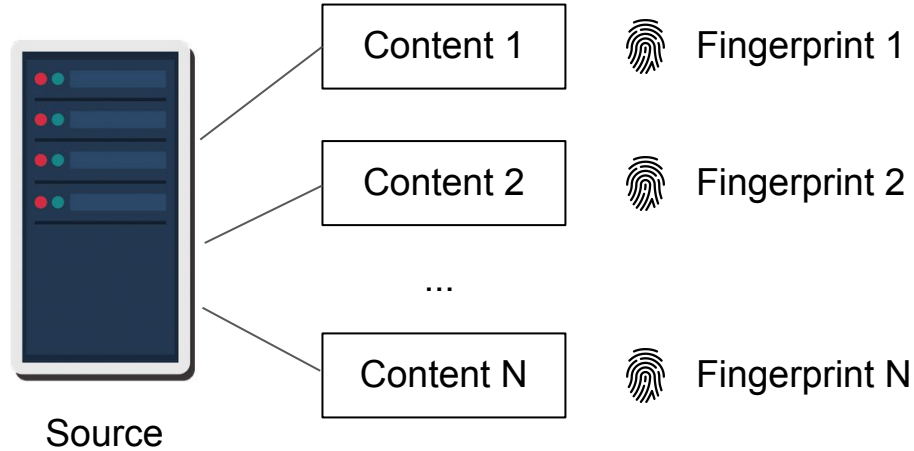
Packet counts

Ordering

etc.

also called a  *Fingerprint*

Often, the source is known



My thesis in one slide

1. Assess the threat, raise awareness



Fingerprint



Something sensitive

2. Present solutions



Reduce, remove or perturb the



Fingerprint

(at a reasonable cost)

Every Byte Matters: Traffic Analysis of Wearable Devices

joint work with Alexandre, Apostolos and Jean-Pierre

Setting

Wearable devices communicate with a smartphone over **Bluetooth**



Fitness tracker



Smartwatch



Sleep tracker



ECG/BPM



The data exchanged is personal and sensitive

Consumer devices

Medical devices



Smartwatch



Fitness tracker



Sleep tracker



ECG/BPM

The data exchanged is personal and sensitive

Consumer devices

Medical devices



Smartwatch



Fitness tracker

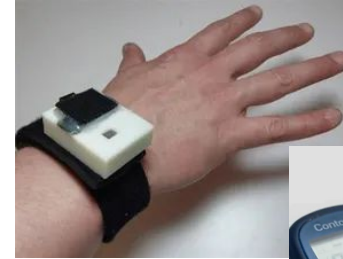


Sleep tracker



ECG/BPM

...



Asthma monitor

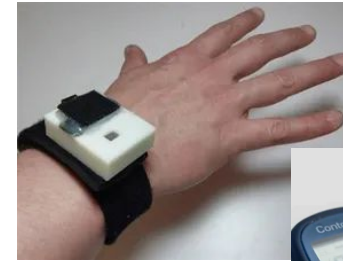


Blood sugar monitor

The data exchanged is personal and sensitive

Consumer devices

Medical devices



- List of apps (app to stop smoking, medication reminder)
- User activities (e.g., recording an insulin injection)

Goal

Do wearable devices produce stable fingerprints,
and what can be inferred then ?

Step 1: buy some devices

Bluetooth Classic

Vendor	Model	OS
Samsung	Galaxy Watch	Tizen
Fossil	Explorist HR	Wear OS 2
Apple	Watch 4	watchOS 5
Huawei	Watch 2	Wear OS 2
Fitbit	Versa 2	Fitbit OS 4
Sony	MDR-XB9	-
Apple	AirPods	-

Bluetooth LE

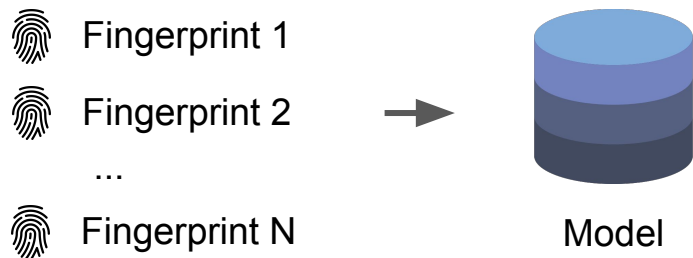
Vendor	Model
Apple	Watch 4
Fitbit	Charge 2
Fitbit	Charge 3
Huawei	Band 3e
Mi	Band 2
Mi	Band 3
Mi	Band 4

- smartwatches
- headphones
- step counters & fitness trackers

Step 2: collect a dataset



Step 3: design & extract fingerprints, train a model



Step 4: identify what can be learned

 unknown fingerprint



Classify

True label \ Predicted label	Add Proteins	Add Calorie	Add Carbs	Add Fat	Add Glucose	Add Insulin
Add Proteins	0.28	0.44	0.03	0.15	0.08	0.02
Add Calorie	0.40	0.46	0.03	0.05	0.05	0.02
Add Carbs	0.01	0.02	0.90	0.03	0.00	0.03
Add Fat	0.07	0.09	0.03	0.81	0.01	0.01
Add Glucose	0.01	0.04	0.00	0.01	0.91	0.03
Add Insulin	0.01	0.01	0.01	0.00	0.01	0.95

(Some) successful attacks:

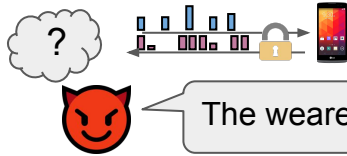
**Device/Action
Identification:**



The wearer is **Measuring the Heart Rate** on its **Huawei Watch 2!**

(Some) successful attacks:

**Device/Action
Identification:**



The wearer is **Measuring the Heart Rate** on its **Huawei Watch 2!**

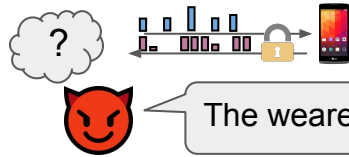
**WearOS Apps
Identification:**



The wearer opened the app **DiabetesM !**

(Some) successful attacks:

**Device/Action
Identification:**



The wearer is **Measuring the Heart Rate** on its **Huawei Watch 2!**

**WearOS Apps
Identification:**



The wearer opened the app **DiabetesM !**

**In-App Action
Identification:**



AddInsulin pressed on application **DiabetesM !**

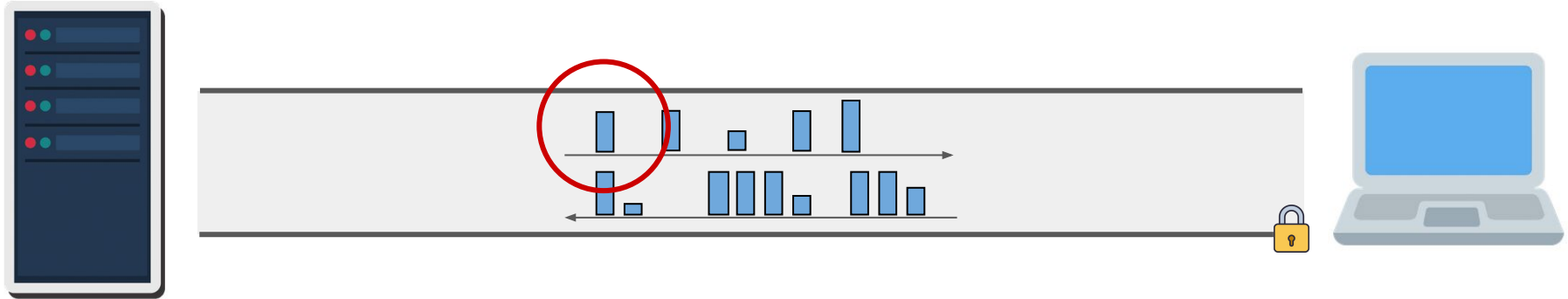


Defenses

- We evaluate 3 standard defenses (pad, delay, and group packets)
- They are **expensive** and **mildly effective** (e.g., 200x data sent, -20% attack accuracy)
- **No easy fix** to the problem, probably **no overarching solution**
- In *specific scenarios*, **defending might be easier**:
 - Wear OS app openings: packet **sizes** matter
 - In-app action fingerprinting: packet **timings** matter
- Other defense strategies: **Data minimization**, **Bulk transfers**

Project 2:

Focus on just one packet. What can we do ?

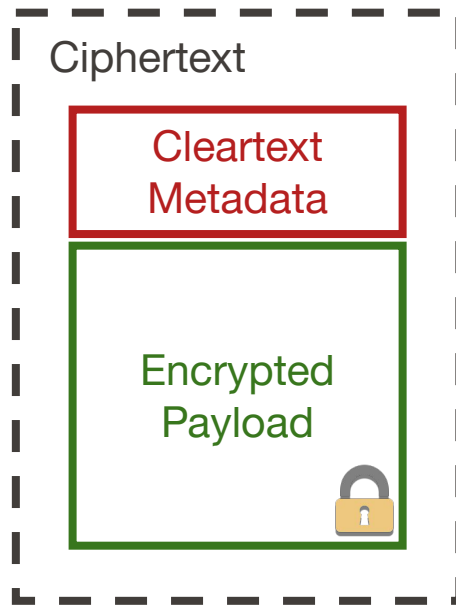


Reducing Metadata Leakage from Static Files

Padmé

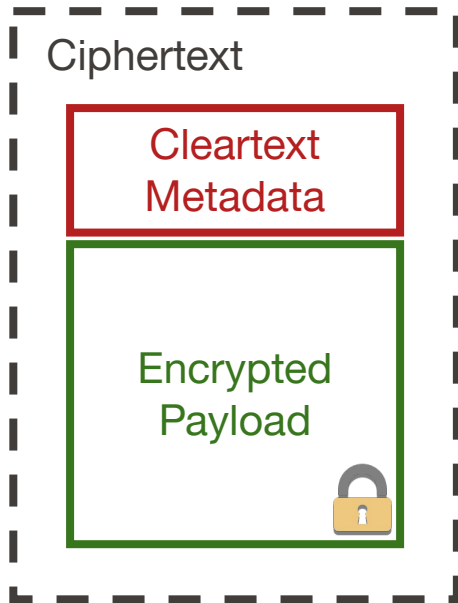
joint work with Kirill, Wouter, Bryan and Jean-Pierre

Ciphertexts often expose metadata in clear



(Contextual Metadata)

Ciphertexts often expose metadata in clear



This **Metadata** may reveal:

What application produced the ciphertext

- it tells the context (could be incriminating in itself)
- it enables searching for implementation errors

The cipher suites used

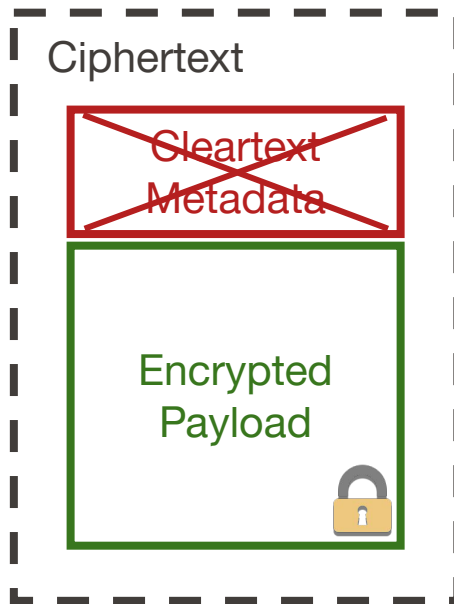
I'm using RC4 !
Crack me !



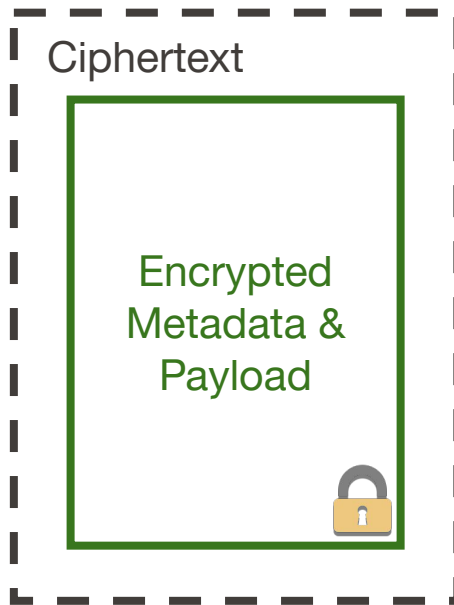
The list of the recipients (PGP)

The length of the contents

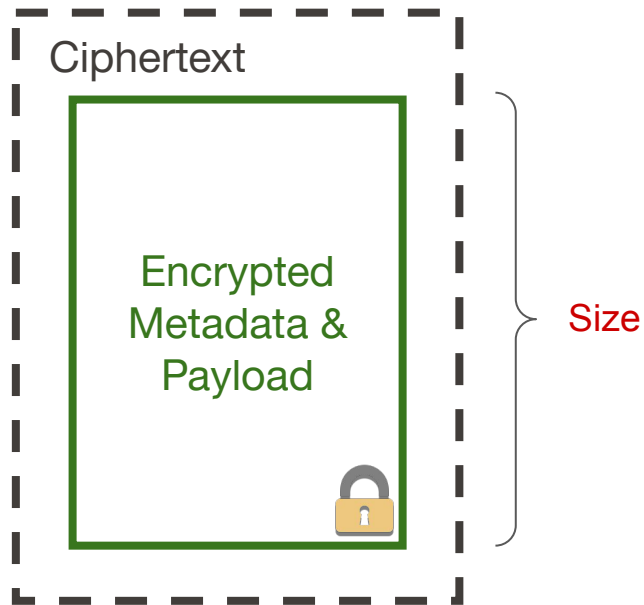
Ciphertexts often expose metadata in clear



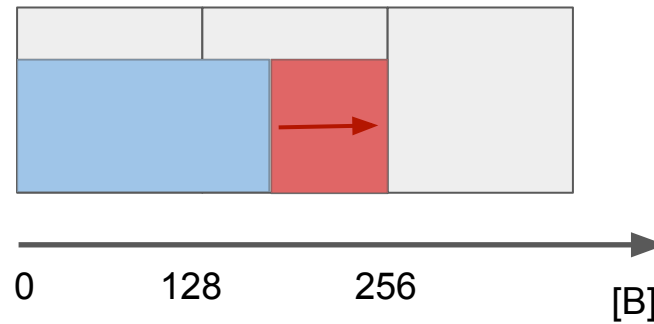
Ciphertexts often expose metadata in clear



Ciphertexts often expose metadata in clear



Constant block-size padding



Constant block-size padding

Problem: no good value for block size

Example: $b = 1 \text{ MB}$



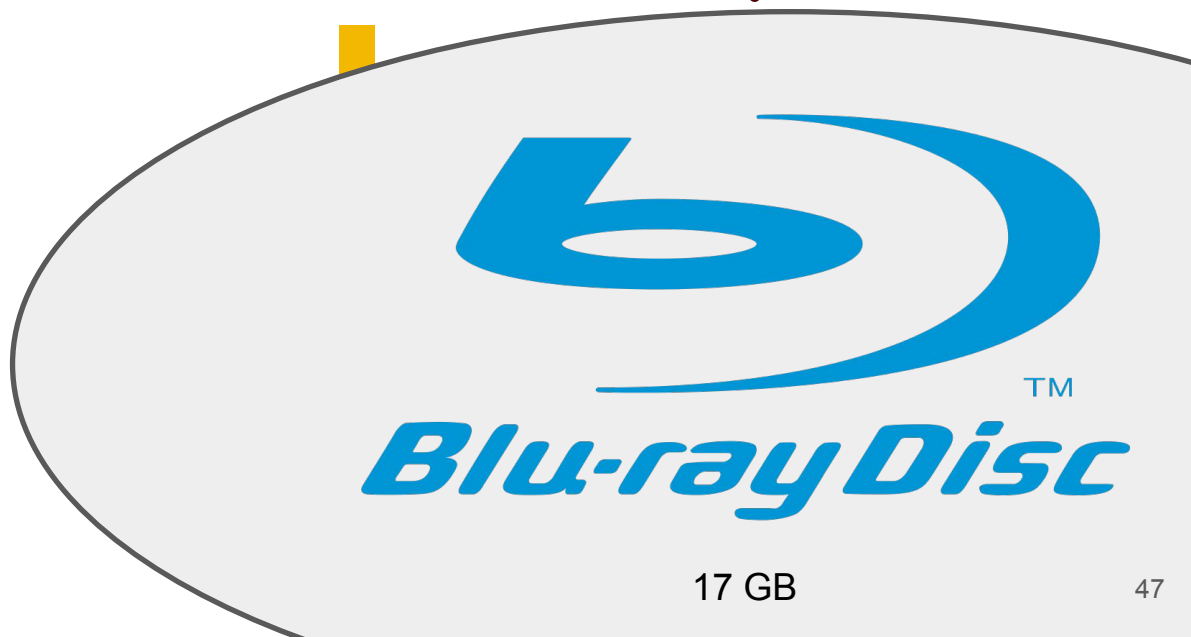
Constant block-size padding

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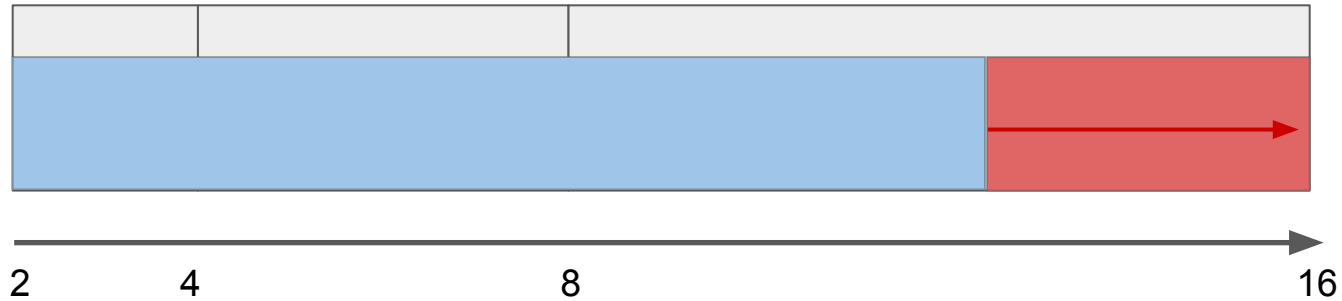


little confusion / privacy



17 GB

Next power of 2



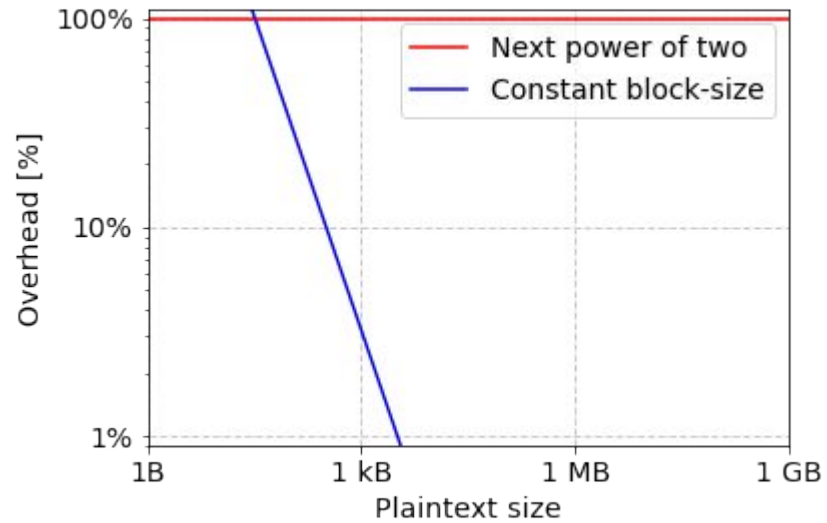
better hiding properties



large overhead (17 GB -> 32 GB)

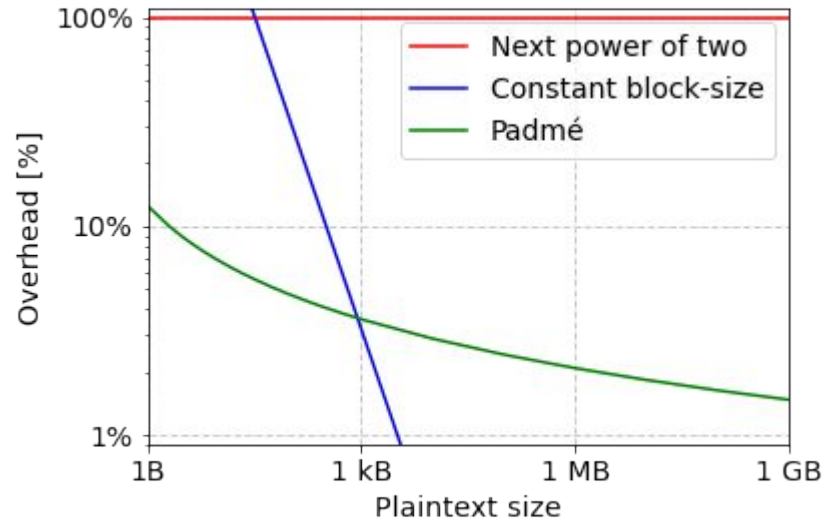
Padmé

Overhead



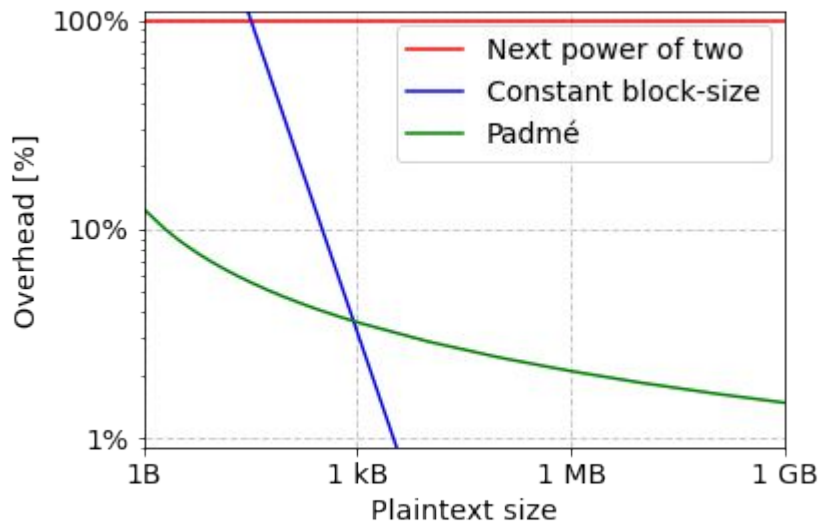
Padmé

Overhead



Padmé

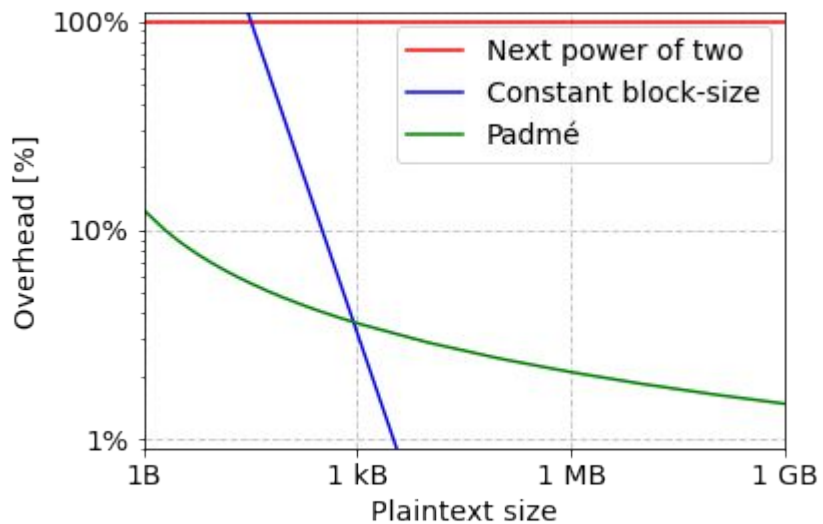
Overhead



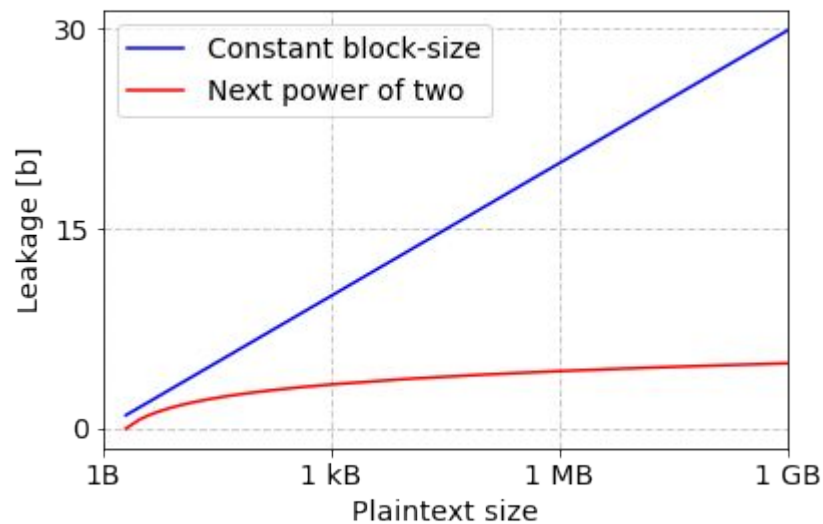
practical overhead

Padmé

Overhead



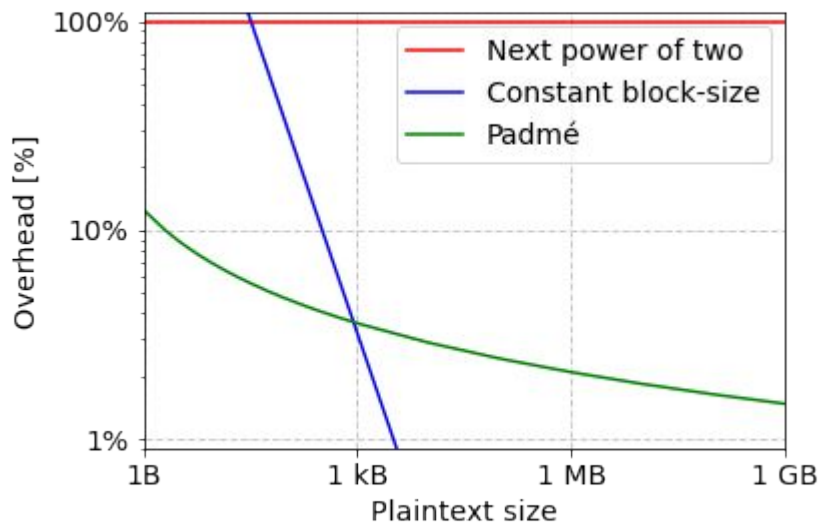
Leakage



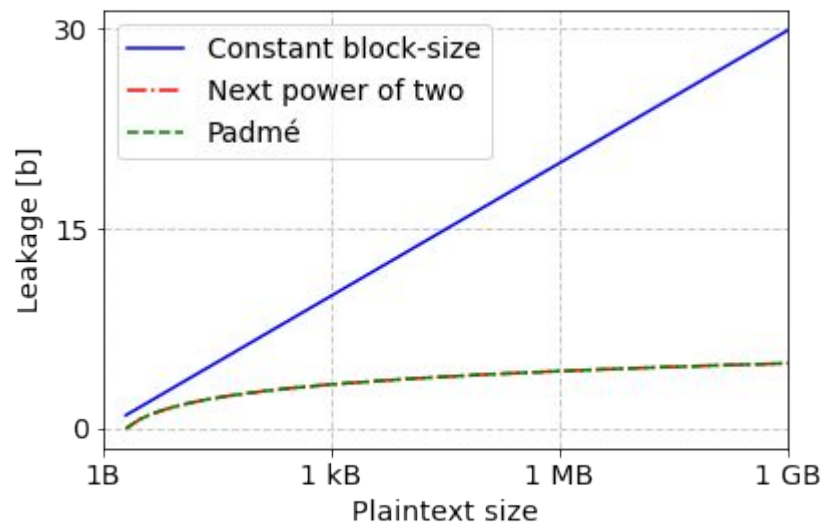
practical overhead

Padmé

Overhead



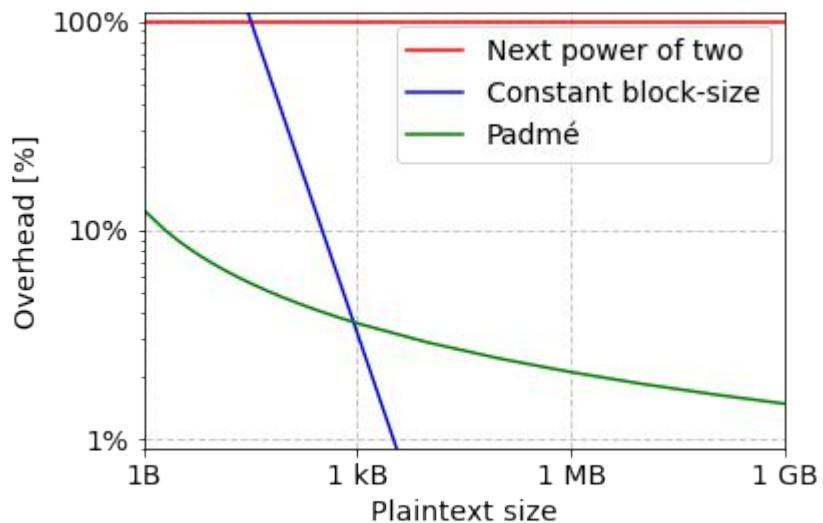
Leakage



practical overhead

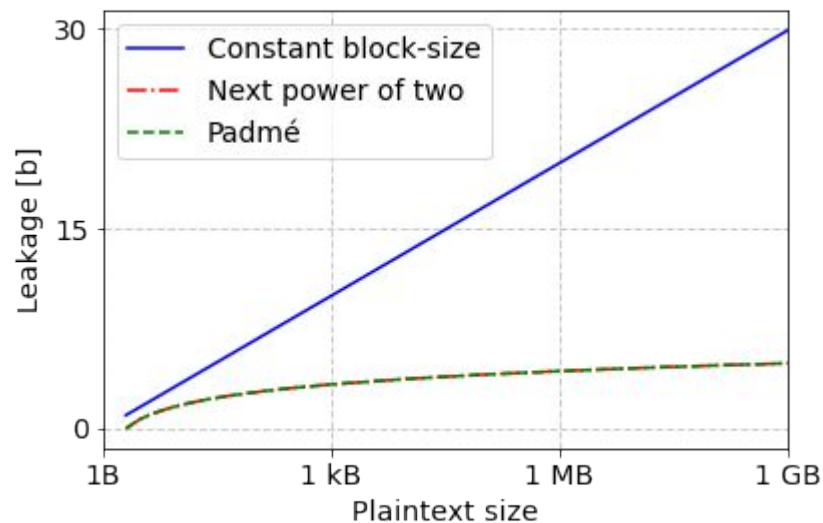
Padmé

Overhead



practical overhead

Leakage



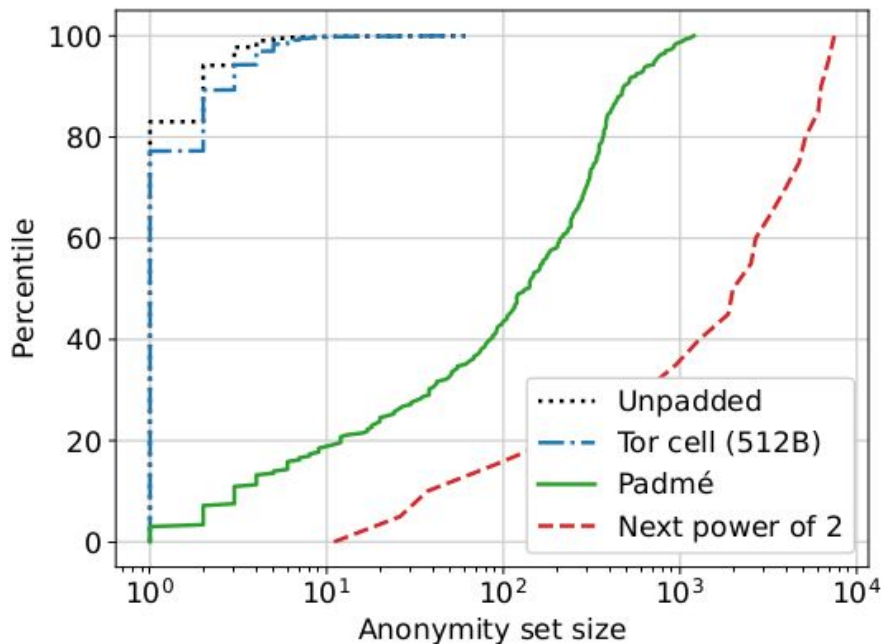
low leakage

Padmé

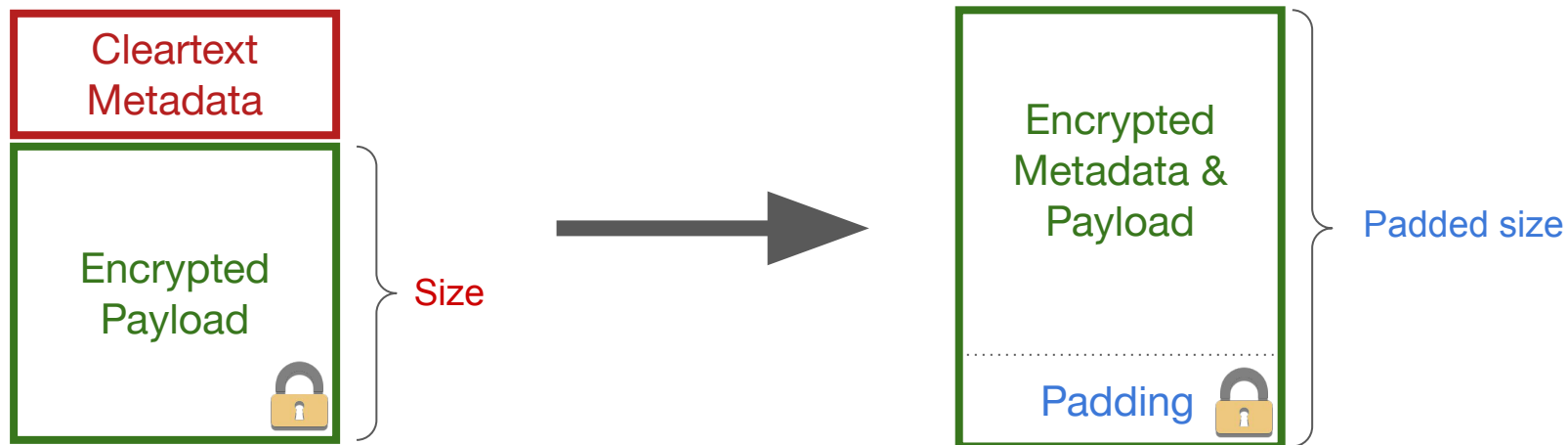
Overhead:

- max +12% $\forall L$
- max +6% $\forall L > 1 \text{ MB}$
- max +3% $\forall L > 1 \text{ GB}$

Ubuntu Packages



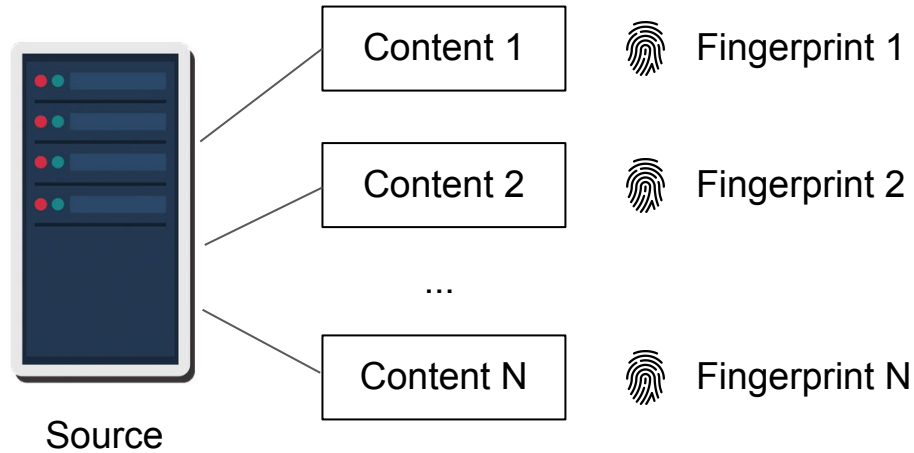
PURBs expose almost no metadata



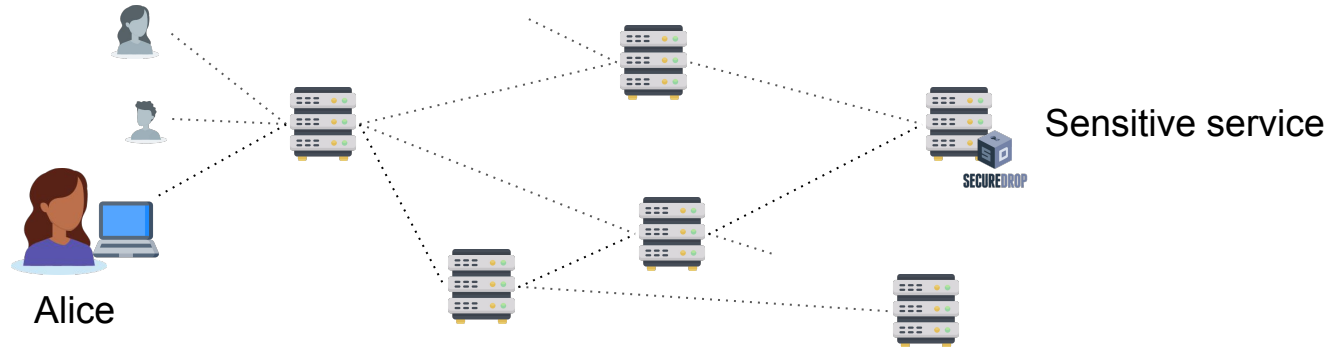
 same functionality

Projects 3 & 4: Anonymous Communication Networks

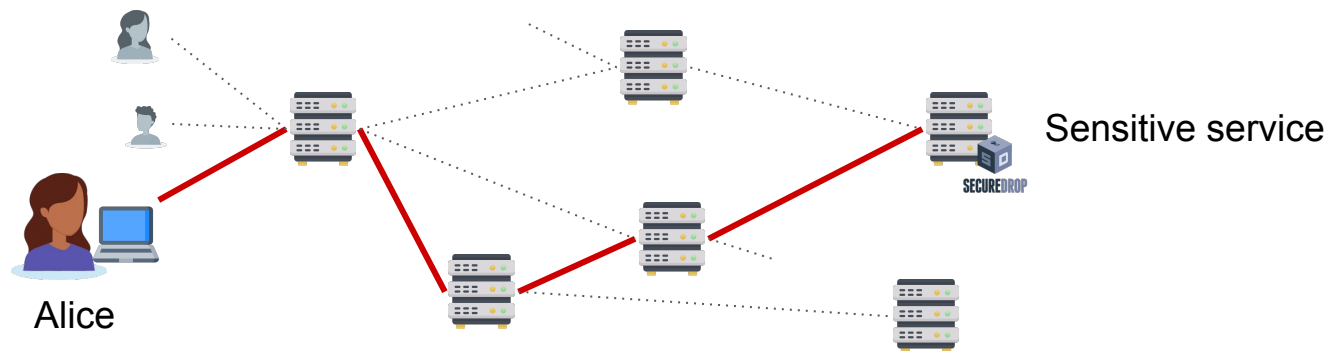
Knowing the source helps a lot



The source / destination can be a sensitive metadata

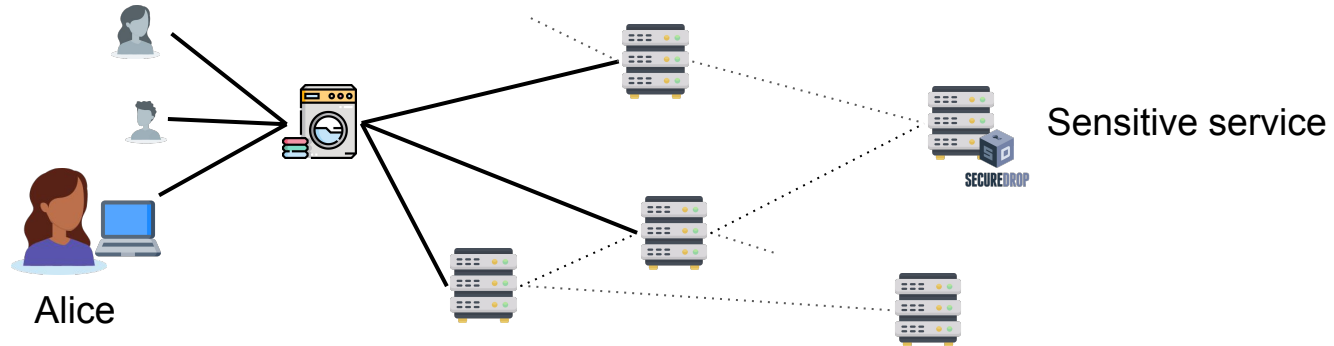


The source / destination can be a sensitive metadata



Alice is contacting SecureDrop !

The source / destination can be a sensitive metadata



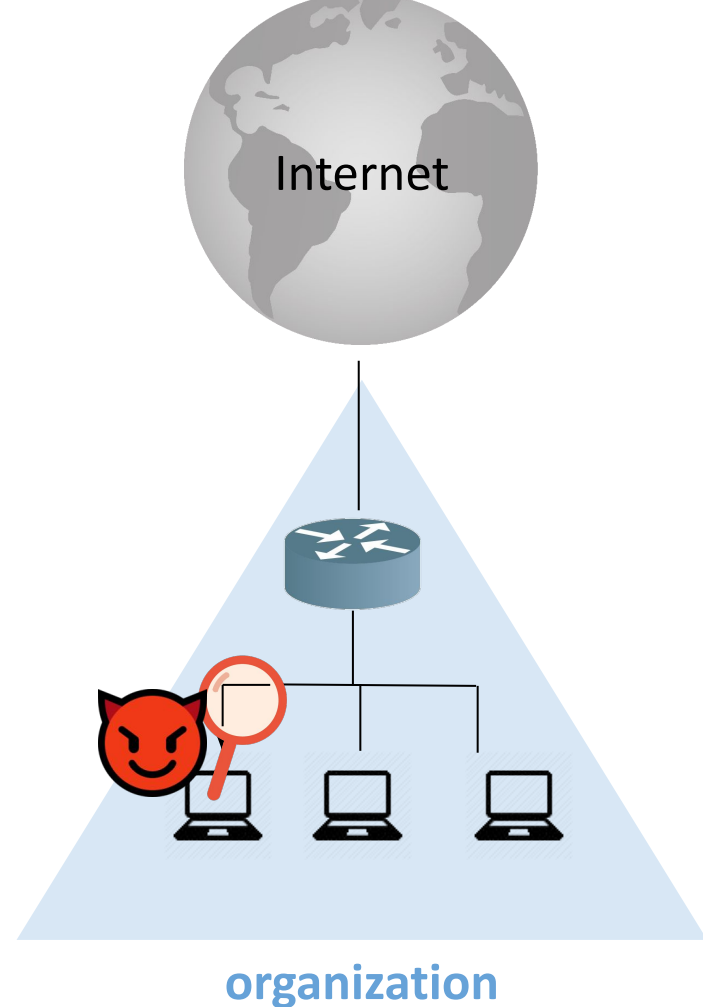
PriFi

A low-latency ACN for LANs and WLANs

joint work with Italo, Apostolos, Bryan and Jean-Pierre

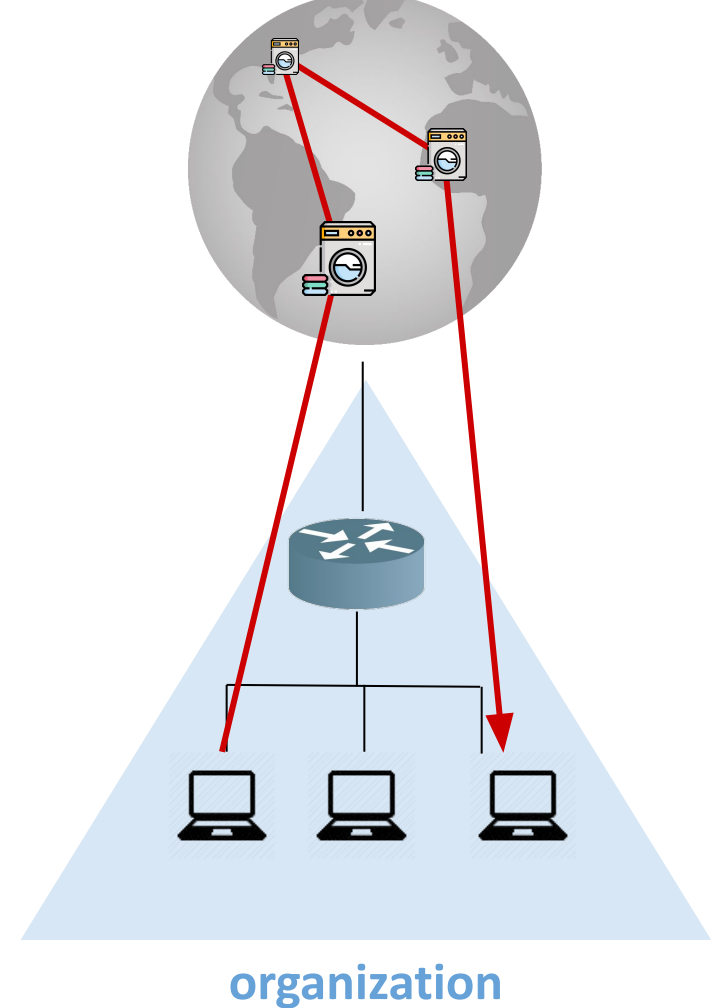
Anonymity for Local-Area Networks

- Problem:
 - Risk of targeted attacks in loosely trusted, sensitive WLANs (e.g., NGOs)
- Goal:
 - Hide source/recipient
 - “Blend in” the traffic of key individuals



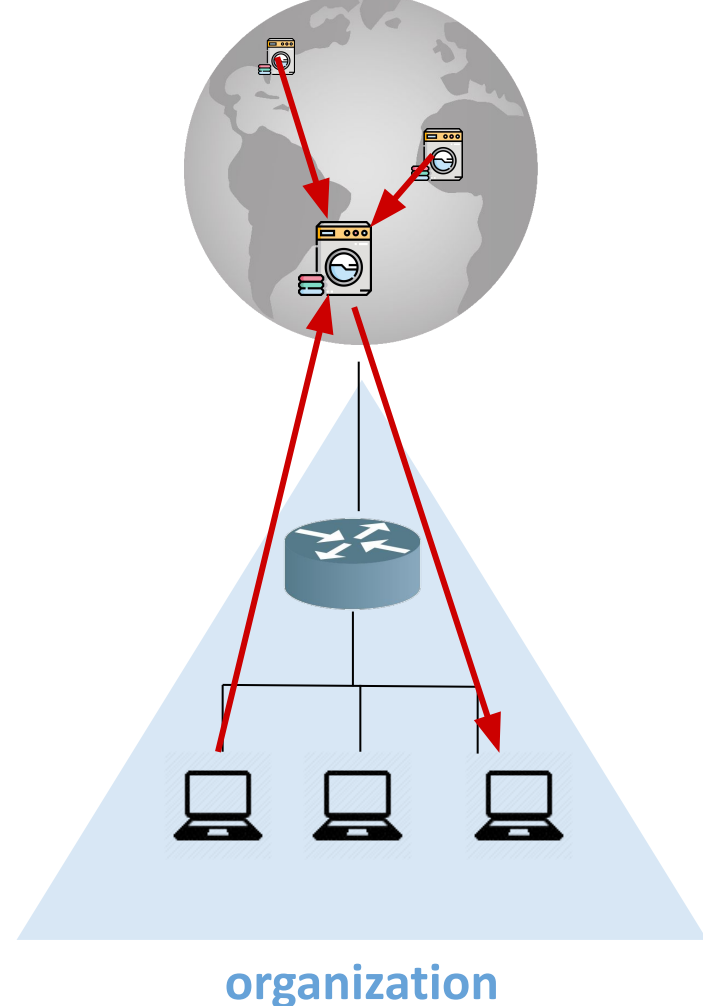
ACNs are poorly suited to LANs

- Tor / Mixnets **add extra hops = extra latency**
- Traffic **leaves the organization**



ACNs are poorly suited to LANs

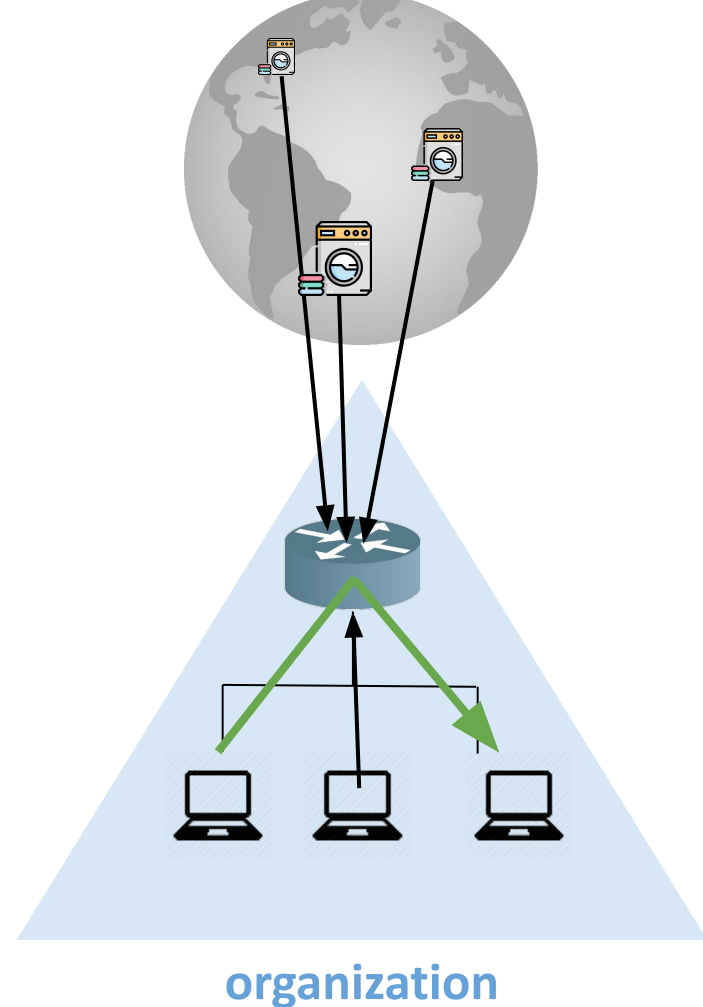
- DC-nets can avoid this
- In practice, **they don't** [10]
- **At each round, “chatty” protocol with the servers** [10]



PriFi

- New topology for DC-nets
- Redesign of the protocols
 - servers contributions are sent in advance
 - avoid server-to-server messages

=> Latency to the servers is not important



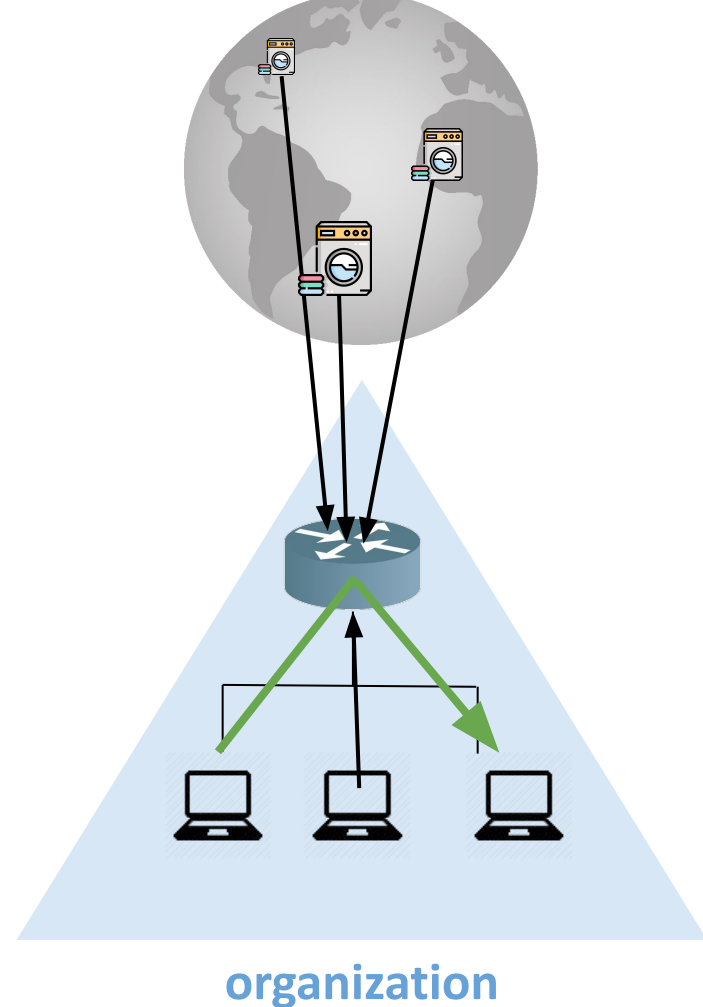
PriFi

- New topology for DC-nets
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=> Latency to the servers is not important

=> “on-path” anonymity

=> cheap broadcast in WLANs



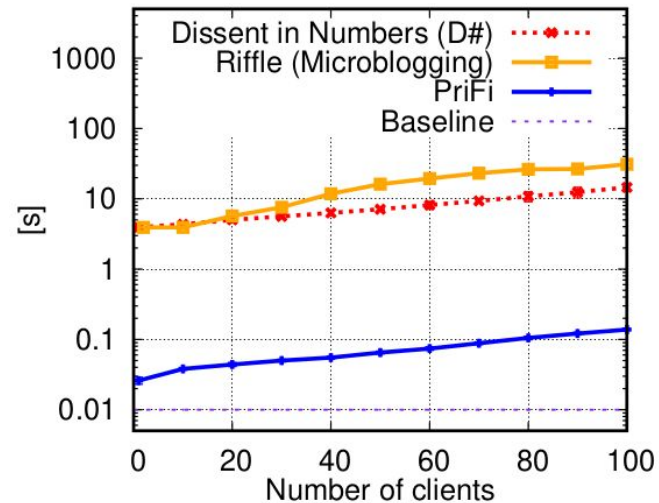
PriFi

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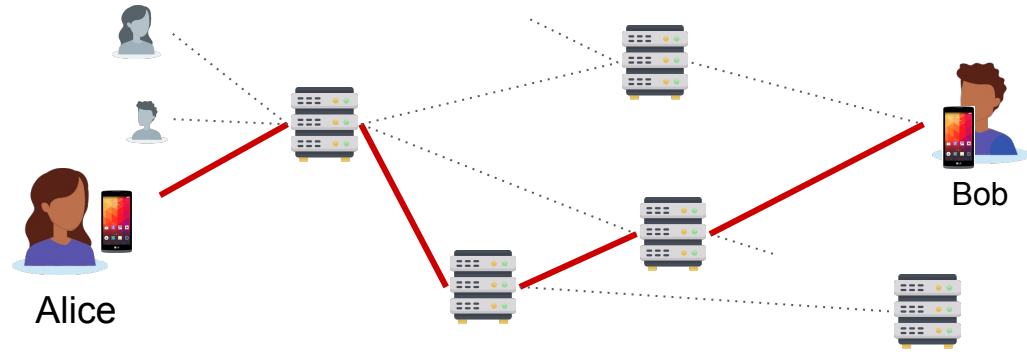
Rubato: Metadata-Private Communications for Mobile Devices

joint work with Moshe, David, Yossi, Nickolai

System for text communication on phones

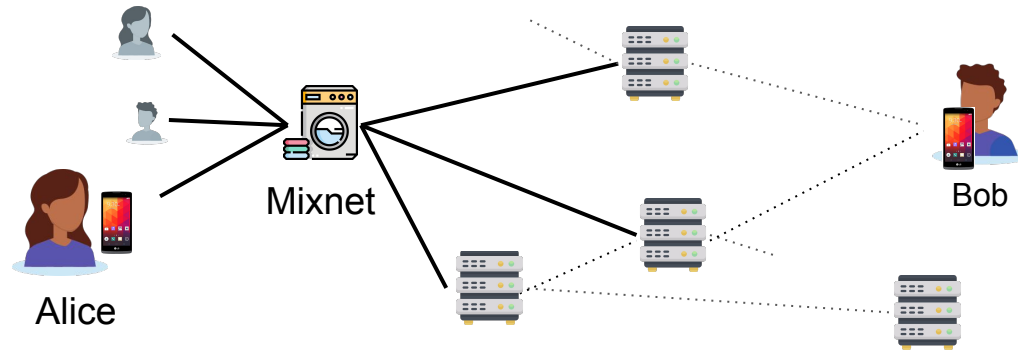


Hide relationships

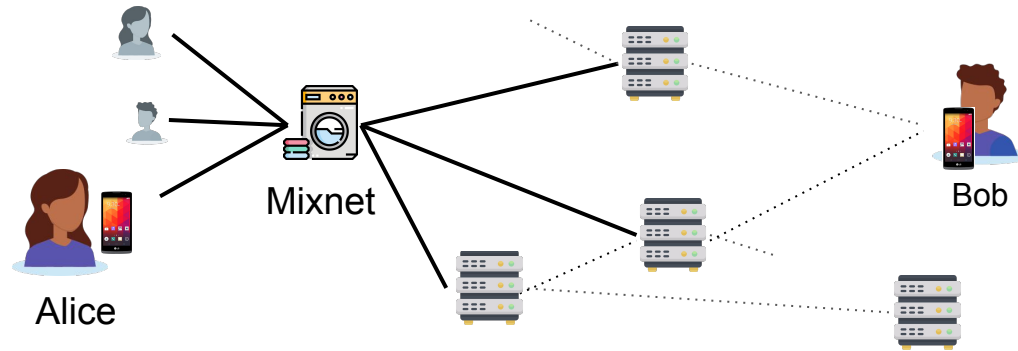


Alice and Bob know each other !

Hide relationships with Mixnets



Hide relationships with Mixnets



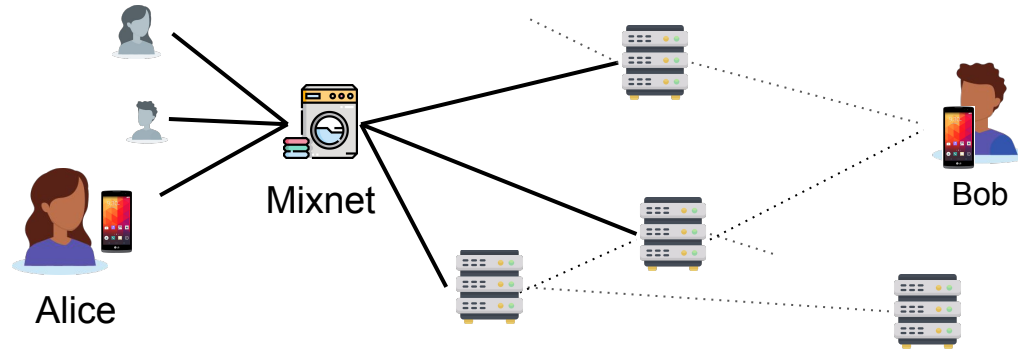
safe



very impractical for users:

- Mixnets supports only one device / user
- it must be always online
- every minute, it must send a message

Hide relationships with Mixnets



safe

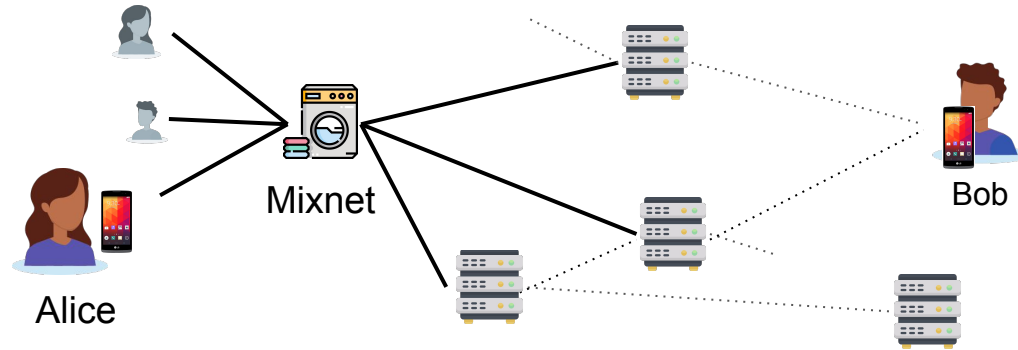


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Hide relationships with Mixnets



safe

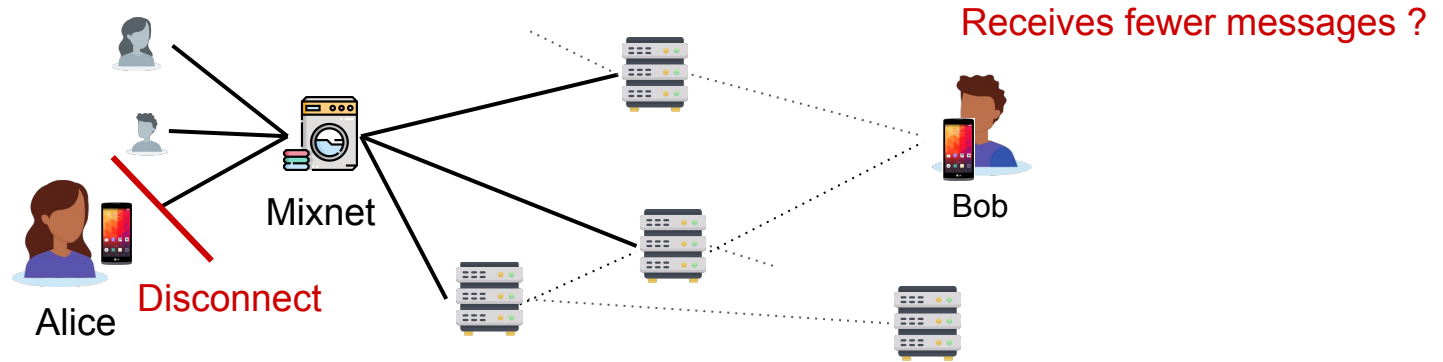


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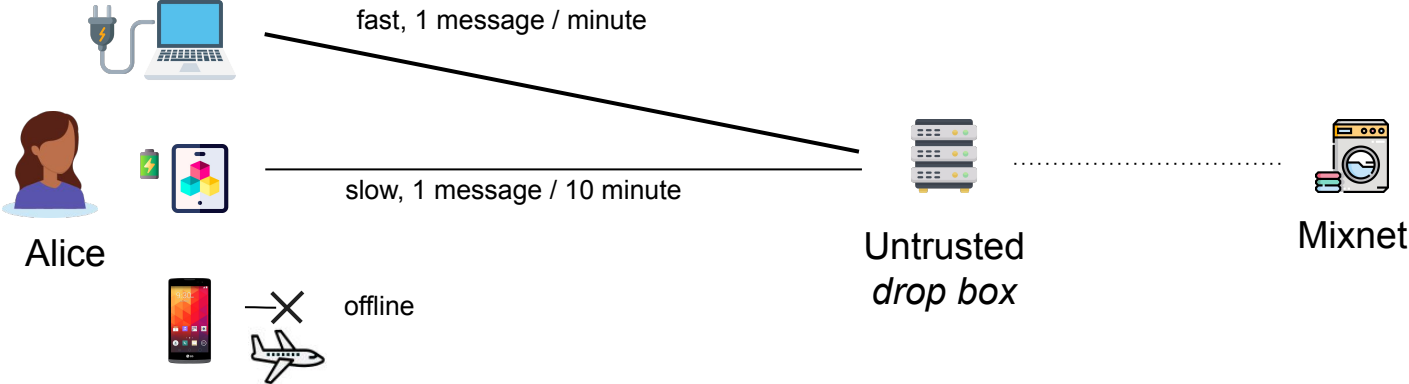
Cost of large-scale mixnets



Global Active Adversary

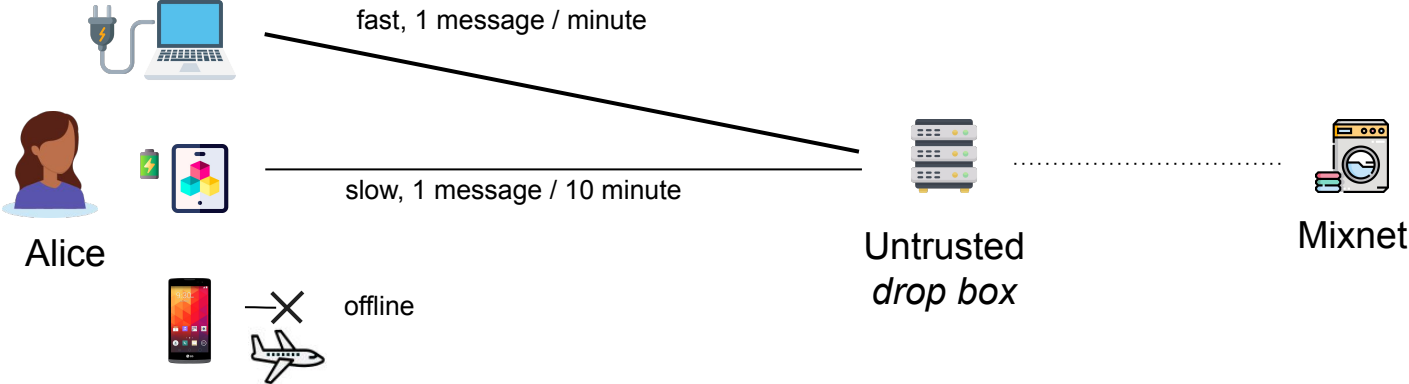
- Sees all the network
- Can delete / delay / inject packets

Rubato



- 👉 user can switch between devices
- 👉 each device can disconnect
- 👉 each device can choose its *schedule*

Rubato



steps Mixnets closer to standard messaging apps



Client Costs

Bandwidth:

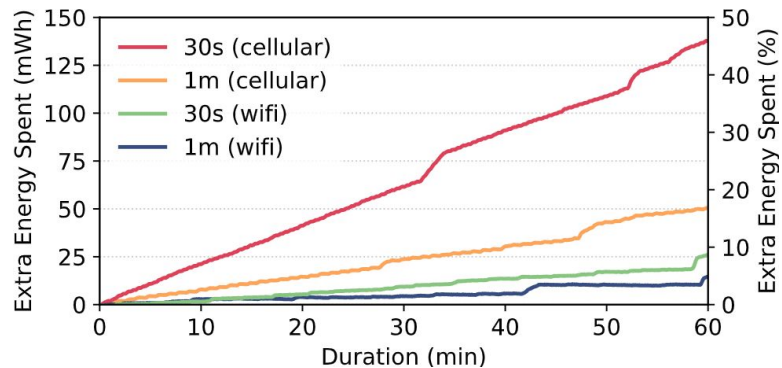
Setup: 110 KB/epoch = 100 MB/month

Messaging: 105 MB/month

Latency: between 32s and 64s

(with a 1-min schedule)

Energy usage:



With a 5-min schedule, after 1h:
≈ +5% energy usage

Conclusion

Contributions of the thesis

- Every Byte Matters: Traffic Analysis of Bluetooth Wearable Devices
 - First **broad analysis of the communication metadata of wearable devices**
 - We **reveal a general susceptibility to traffic-analysis attacks**, which can allow:
 - **identifying devices, applications, user actions**
 - **tracking and profiling users**
 - If we want to protect such information, we need defense strategies

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 - If we want to protect such information, we need defense strategies
- Padmé
 - Padding function with **low costs (<12%)** that **outperforms classic approaches** asymptotically
 - In practice, we show that it has **good hiding properties**

Contributions of the thesis (cont')

- PriFi
 - Low-latency, traffic-agnostic anonymity for a small set of users (VoIP support)
 - The latency does not depend on the latency to the anytrust servers
 - “On-path” anonymization that provides low latency

Contributions of the thesis (cont')

- PriFi
 - Low-latency, traffic-agnostic anonymity for a small set of users (VoIP support)
 - The latency does not depend on the latency to the anytrust servers
 - “On-path” anonymization that provides low latency
- Rubato
 - First large-scale ACN with multi-device, asynchronous clients (Global Active Adversary setting)
 - Each device can choose its communication frequency & costs
 - It enables mobile devices to participate at a reasonable cost

Impact outside of research

- Every Byte Matters: Traffic Analysis of Bluetooth Wearable Devices
 - Contacted ~100 vendors and manufacturers, ~10 follow-ups by email, 2 follow-up meetings with large device manufacturers
 - Received a bug bounty
- Padmé
 - Maintainers of SequoiaPGP implemented Padmé
- PriFi
 - Demos at the Red Cross (ICRC) headquarters and at EPFL (one awarded a prize)
 - Patent

Next steps for metadata privacy ?

It is still an open problem

- No one-size-fits-all defense
 - => Per domain, iteratively evaluate risks
- Compared to non-metadata-private alternatives, solutions are costly
 - => Increase visibility of the attacks to justify the costs

Building safer apps

- Could we have (automated?) guidelines for app developers ?

- Could we have “defense strategies” provided by the OS ?

This could be an opportunity for **designing the defenses iteratively**



Analyzing and Protecting Communication Metadata

Experimental

Every Byte Matters [4]

Traffic-analysis attack of wearable devices.

Theoretical

Padmé [2]

A padding function that efficiently hides sizes.

Attacks

Defenses

Systems: Anonymous Communication Networks (ACN)

PriFi [1,3]

Traffic-agnostic, low-latency ACN for local-area networks.

Rubato [5]

Large-scale ACN for text messaging on mobile devices.

Thesis & technical presentations soon on <https://lbarman.ch>

[1] L. Barman, M. Zamani, I. Dacosta, J. Feigenbaum, B. Ford, J.-P. Hubaux, D. Wolinsky. **PriFi: A Low-latency [...] Protocol for Local-Area Anonymous [...]**. WPES 2016.

[2] K. Nikitin*, L. Barman*, W. Lueks, M. Underwood, J.-P. Hubaux, B. Ford. **Reducing Metadata Leakage from Encrypted Files and Communication with PURBs**. PETS 2019

[3] L. Barman, I. Dacosta, M. Zamani, E. Zhai, A. Pyrgelis, B. Ford, J. Feigenbaum, J.-P. Hubaux. **PriFi: Low-latency Anonymity for Organizational Networks**. PETS 2020

[4] L. Barman, A. Dumur, A. Pyrgelis, J.-P. Hubaux. **Every Byte Matters: Traffic Analysis of Bluetooth Wearable Devices**. UbiComp 2021.

[5] L. Barman, M. Kol, D. Lazar, Y. Gilad, N. Zeldovich. **Rubato: Metadata-Private Messaging for Mobile Devices**. Under submission.